

LECTURE 3:
MIXED STRATEGIES

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Games and Decisions

Matrix Games (revision)

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- a special case of zero-sum games:
 - a finite set of agents: $\{1,2\}$
 - strategy spaces (*finite*): $\{X,Y\}$
 - strategy profile: (x,y)
 - payoff functions: $Z_1(x,y), Z_2(x,y)$
 - zero-sum payoffs: $Z_1(x,y) + Z_2(x,y) = 0$

- payoffs written in a matrix, typically denoted by A :

$$A = (a_{ij})_{\substack{i=1,\dots,m \\ j=1,\dots,n}} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix}$$

- $a_{ij} =$ the payoff of player 1 for strategy profile (i,j)
(i.e., player 1 picks i th strategy and player 2 picks j th)

Nash Equilibrium in Matrix Games (revision)

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□ mathematical definition:

A strategy profile (x^, y^*) with the property that*

$$Z_1(x, y^*) \leq Z_1(x^*, y^*) \leq Z_1(x^*, y)$$

for all $x \in X$ and $y \in Y$ is a NE.

Inequality from the definition above explained:

→ If player 1 deviates from NE, he/she won't be any better off

$$Z_1(x, y^*) \leq Z_1(x^*, y^*) \leq Z_1(x^*, y)$$

← If player 2 deviates from NE, he/she won't be any better off

Matrix Games with No Pure Strategy NE

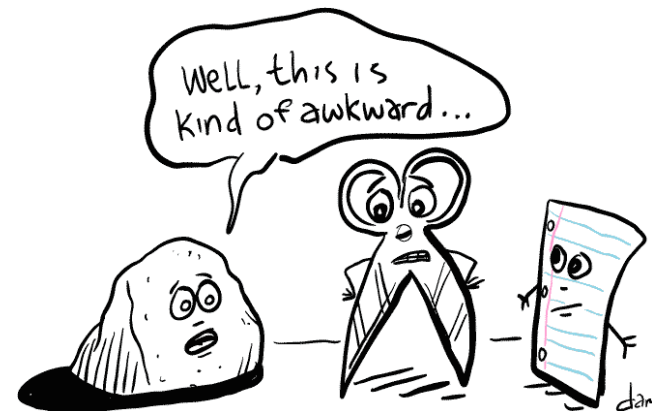
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Rock, Paper, Scissors

Player 2

1 \ 2	R	P	S
R	0	-1	1
P	1	0	-1
S	-1	1	0

Player 1



Matrix Games with No Pure Strategy NE (cont'd)

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Matching pennies game

- both players secretly turn a coin to heads/tails
- if both heads or both tails, player 1 pays \$1, otherwise player 2 pays \$1

		Player 2	
		Heads	Tails
Player 1	Heads	-1	1
	Tails	1	-1



Matrix Games with No Pure Strategy NE (cont'd)

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Penalty kicks (a modified version of matching pennies)

- Kick vs. Goalkeeper
- *strategies*: Right/Left
- *payoffs*: scoring probabilities

		Goalkeeper	
		Right	Left
Kick	Right	0.5	0.9
	Left	0.9	0.5



Matrix Games with No Pure Strategy NE (cont'd)

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- *stability* in NE: even if he knows his opponent is playing the NE strategy, player 1 has no incentive to deviate from NE (the same goes for player 2)
- matching pennies:
 - if player one knows player 2 is playing Heads, he can win by playing Tails and vice versa
 - neither Heads nor Tails can be the NE strategy
- still, there is a rational way to play the game:
 - each time you play the game, toss the coin and let it decide the strategy on itself
 - even if your opponent knows your strategy, she can't take advantage of that (compare with *RPS* game)
 - tossing a coin means applying a *mixed strategy*

Pure Vs. Mixed Strategies in Matrix Games

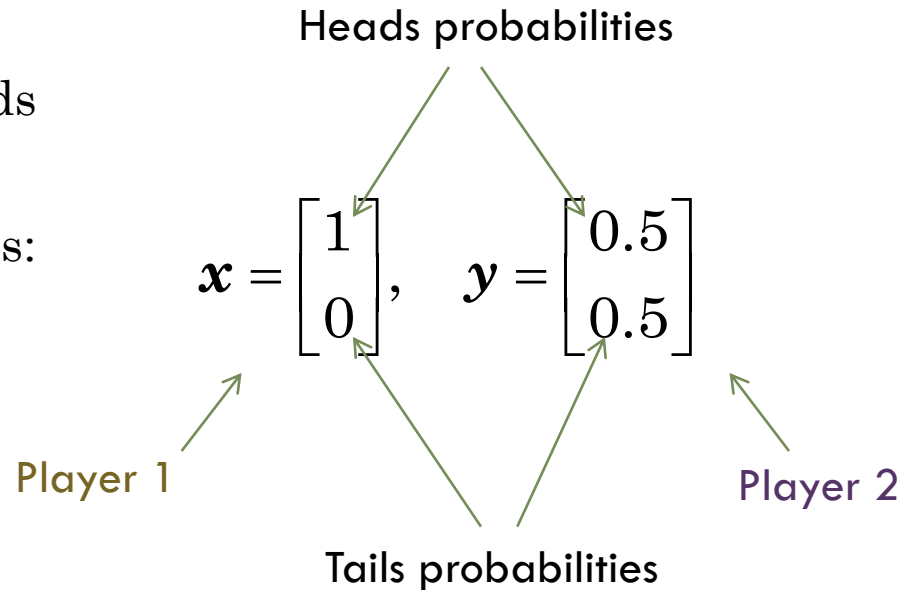
- switch from *pure strategies* to *mixed strategies*:
 - *pure strategy*: the player decides for a particular strategy (i.e., picks a certain row or column)
 - *mixed strategy*:
 - the player decides about the probabilities of the alternative strategies
(*sum of the probabilities = 1*)
 - when the decisive moment comes, he/she makes a random selection of the strategy with the stated probabilities
- allowing for mixed strategies in a matrix game: sometimes called a *mixed extension of a matrix game*
- *notation*: mixed strategies = column vectors \mathbf{x} and \mathbf{y} , *i*th element represents the probability of *i*th row/column being picked

Mixed Strategies: An Example

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- matching pennies game
 - player 1 always chooses Heads
 - player 2 tosses a coin
- the resulting mixed strategies:

		Player 2	
		Heads	Tails
Player 1	Heads	-1	1
	Tails	1	-1



- *note: \mathbf{x} actually is a pure strategy here*

Payoffs with Mixed Strategies

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- when using mixed strategies, the payoffs for the individual players become *random variables*
 - the possible values are stated in the game's matrix
 - the final outcome depends on the strategies eventually picked by the individual players
- in order to be able to treat the payoffs resulting from a combination of mixed strategies as a single number, we use the concept of *expected payoff*

Payoffs with Mixed Strategies

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- when using mixed strategies, the payoffs for the individual players become *random variables*
 - ▣ the possible values are stated in the game's matrix
 - ▣ the final outcome depends on the strategies eventually picked by the individual players
- in order to be able to treat the payoffs resulting from a combination of mixed strategies as a single number, we use the concept of *expected payoff*

Expected value:

Let Z be a discrete-valued random variable with possible values in Ω . The *expected value of Z* is:

$$EZ = \sum_{z \in \Omega} z \cdot \Pr\{Z = z\}$$

Expected Payoffs: An Example

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- consider matching pennies again, with $\mathbf{x} = [1 \ 0]^T$, $\mathbf{y} = [0.5 \ 0.5]^T$
 - the expected payoff of player one is:

$$\begin{aligned} \mathbf{E}Z_1 &= \begin{Bmatrix} -1 \times \Pr(H, H) & +1 \times \Pr(H, T) \\ +1 \times \Pr(T, H) & -1 \times \Pr(T, T) \end{Bmatrix} \\ &= \begin{Bmatrix} -1 \times (1 \times 0.5) & +1 \times (1 \times 0.5) \\ +1 \times (0 \times 0.5) & -1 \times (0 \times 0.5) \end{Bmatrix} \\ &= \begin{Bmatrix} -0.5 & +0.5 \\ +0 & +0 \end{Bmatrix} = 0 \end{aligned}$$

		0.5	0.5
	1 \ 2	Heads	Tails
1	Heads	-1	1
0	Tails	1	-1

- the expected value of player 2 is:

$$\mathbf{E}Z_2 = -\mathbf{E}Z_1 = 0$$

Expected Payoffs: A Generalization

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- consider a general 2×2 matrix game with matrix \mathbf{A} and mixed strategies \mathbf{x} and \mathbf{y} :

$$\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}, \quad \mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}, \quad \mathbf{A} = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}.$$

	y_1	y_2
x_1	a_{11}	a_{12}
x_2	a_{21}	a_{22}

- the expected payoff of player 1:

$$\begin{aligned} \mathbb{E}Z_1 &= x_1 a_{11} y_1 + x_1 a_{12} y_2 \\ &\quad + x_2 a_{21} y_1 + x_2 a_{22} y_2 = \sum_i \sum_j x_i a_{ij} y_j = \mathbf{x}^\top \mathbf{A} \mathbf{y} \end{aligned}$$

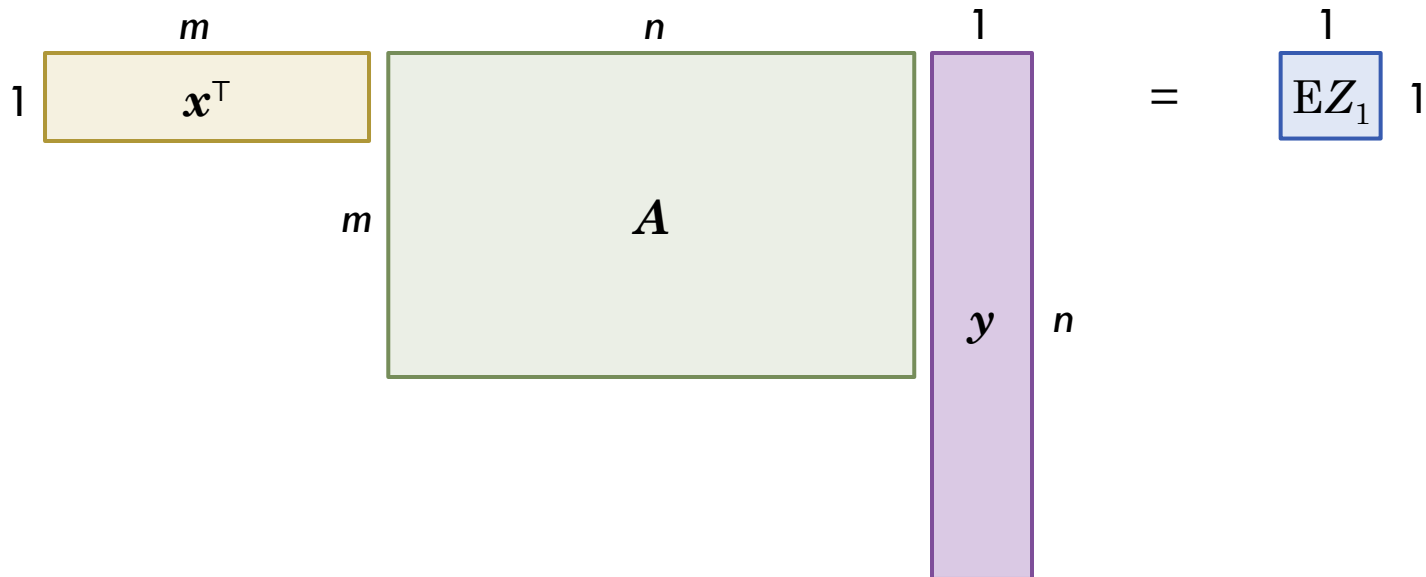
- the last two expressions hold for the general case of $m \times n$ matrix games

Expected Payoffs: A Generalization (cont'd)

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Expected payoffs in an $m \times n$ matrix game with mixed strategies:

$$EZ_1 = \sum_{i=1}^m \sum_{j=1}^n x_i a_{ij} y_j = \mathbf{x}^\top \mathbf{A} \mathbf{y} = -EZ_2$$



Expected payoffs: Exercise 1

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- calculate the expected payoff for both players in the penalty kicks game for the following mixed strategies of the two players:
 - a) $\mathbf{x} = [1 \ 0]^T$, $\mathbf{y} = [0.5 \ 0.5]^T$.
 - b) $\mathbf{x} = [0.5 \ 0.5]^T$, $\mathbf{y} = [0.5 \ 0.5]^T$.
 - c) $\mathbf{x} = [0 \ 1]^T$, $\mathbf{y} = [0.5 \ 0.5]^T$.

		Goalkeeper	
		Right	Left
Kick	1 \ 2 Right	0.5	0.9
	Left	0.9	0.5

Nash Equilibrium with Mixed Strategies

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- overall concept: the same as with pure strategies

“NE is such a combination of strategies that neither of the players can increase their payoff by choosing a different strategy.”

“NE is a solution with the property that whoever of the players chooses some other strategy, he or she will not increase his or her payoff.”

- **mathematical definition:**

NE is a combination of (mixed) strategies \mathbf{x}^ and \mathbf{y}^* with the property that*

$$\mathbf{x}^\top \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}$$

for all mixed strategies \mathbf{x} and \mathbf{y} .

- **value of the game:** player 1's expected payoff at NE ($\mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^*$)

Nash Equilibrium with Mixed Strategies

(cont'd)

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Inequality from the NE definition explained:

Can be rewritten as:

$$EZ_1(\mathbf{x}, \mathbf{y}^*) \leq EZ_1(\mathbf{x}^*, \mathbf{y}^*),$$

which means: If player 1 deviates from NE, his/her expected payoff will not increase

$$\mathbf{x}^\top \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}$$

Can be rewritten as:

$$-EZ_2(\mathbf{x}^*, \mathbf{y}^*) \leq -EZ_2(\mathbf{x}^*, \mathbf{y}), \text{ or } EZ_2(\mathbf{x}^*, \mathbf{y}) \leq EZ_2(\mathbf{x}^*, \mathbf{y}^*),$$

which means: If player 2 deviates from NE, his/her expected payoff will not increase

Basic Theorem on Matrix Games

- it is easily seen that a pure-strategy NE is also a mixed-strategy NE; therefore, if a matrix game has a NE in pure strategies, it has a NE in mixed strategies as well
- what happens in the are no pure-strategy NE? As JOHN VON NEUMANN proved in 1928, even if a matrix game has no NE in pure strategies (i.e., no saddle point of the payoff matrix), it still has a NE in mixed strategies (*always*)

- **Basic Theorem on Matrix Games:**

For any matrix A there exist mixed strategies \mathbf{x}^ and \mathbf{y}^* such that*

$$\mathbf{x}^\top \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^* \leq \mathbf{x}^{*\top} \mathbf{A} \mathbf{y}$$

for all mixed strategies \mathbf{x} and \mathbf{y} .

Finding Mixed-Strategy NE's

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- Basic Theorem tells us that every matrix game has a mixed-strategy solution
- two possible approaches of finding the solution depending on the size of the matrix A :
 - a) “small” – $2 \times n$ and $m \times 2$ matrices:
 - simple
 - *graphical solution*
 - b) “general” – $m \times n$ matrices:
 - a bit more complicated
 - *linear programming*
 - can be solved in *MS Excel* (or other SW)

Graphical Solution – 2×2 Matrices

(cont'd)

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- 2×2 matrix \rightarrow players actually choose only 1 probability:

$$x_2 = 1 - x_1, \quad y_2 = 1 - y_1$$

- player 1's decision: x_1 (probability of T)

- expected payoffs:

- player 2 picks L : $EZ_1 = 1 \times x_1 + 11 \times (1 - x_1) = 11 - 10x_1$

- player 2 picks R : $EZ_1 = 7 \times x_1 + 2 \times (1 - x_1) = 2 + 5x_1$

- player 2 chooses a mixed strategy: anything in between (more precisely: a *convex combination* of the two)

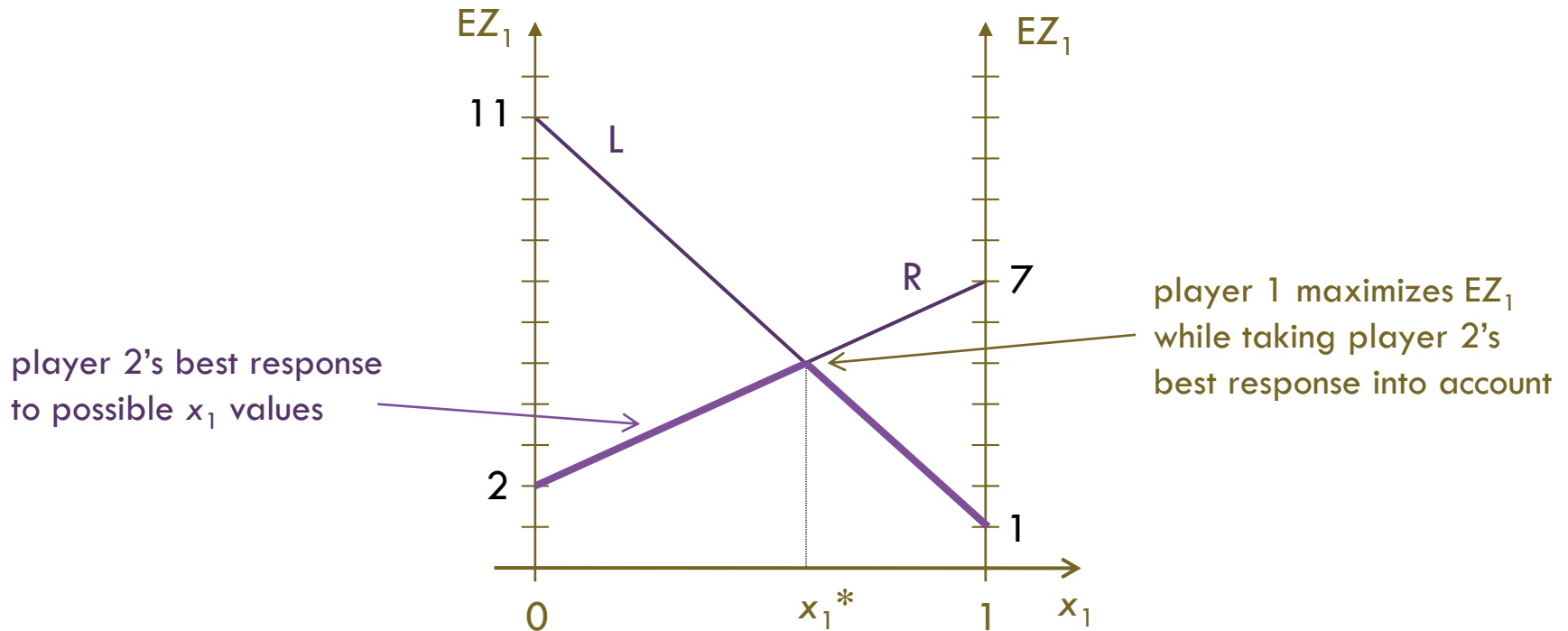
		Player 2	
	1 \ 2	L	R
Player 1	T	1	7
	B	11	2

	y_1	$1 - y_1$
x_1	a_{11}	a_{12}
$1 - x_1$	a_{21}	a_{22}

Graphical Solution – 2×2 Matrices

(cont'd)

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Graphical Solution – 2×2 Matrices

(cont'd)

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- player 2's decision: y_1 (probability of L)
 - expected payoffs:
 - player 1 picks T : $EZ_1 = 1 \times y_1 + 7 \times (1 - y_1) = 7 - 6y_1$
 - player 1 picks B : $EZ_1 = 11 \times y_1 + 2 \times (1 - y_1) = 2 + 9y_1$
 - player 1 chooses a mixed strategy: anything in between (more precisely: a *convex combination* of the two)

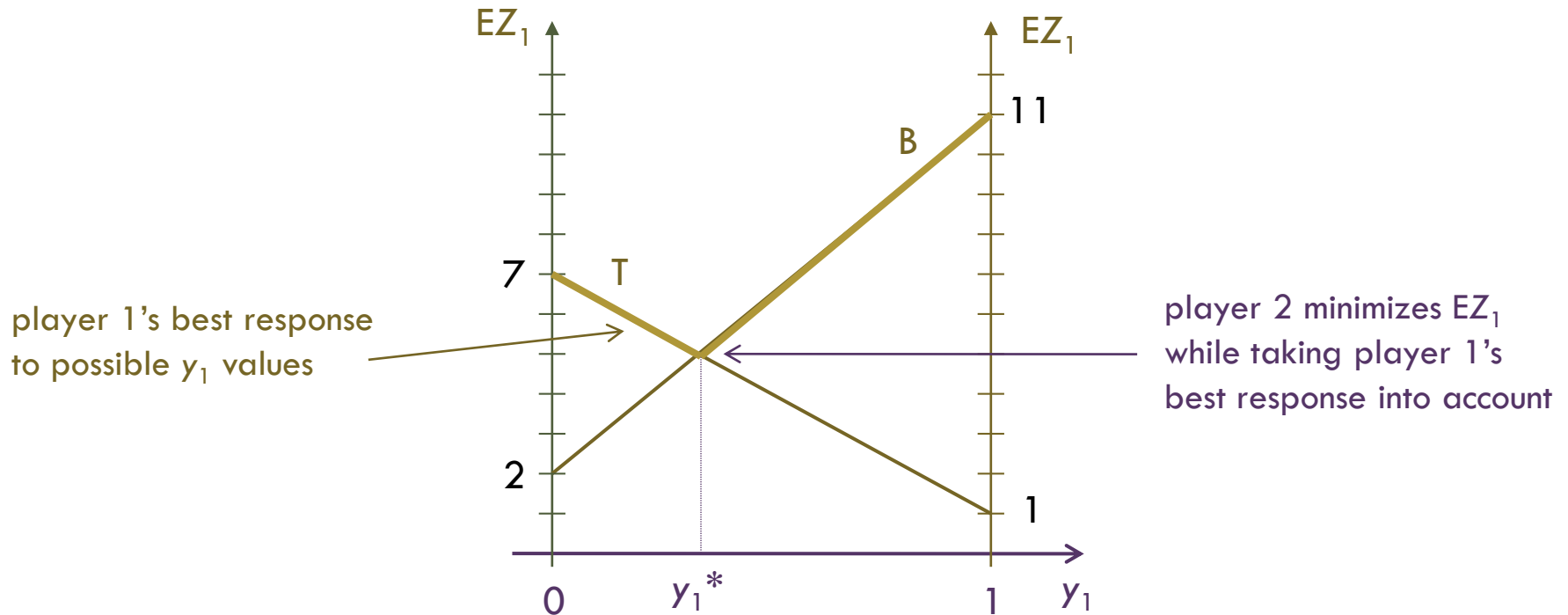
		Player 2	
		L	R
Player 1	T	1	7
	B	11	2

	y_1	$1 - y_1$
x_1	a_{11}	a_{12}
$1 - x_1$	a_{21}	a_{22}

Graphical Solution – 2×2 Matrices

(cont'd)

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Graphical Solution – 2×2 Matrices

(cont'd)

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- finding x_1^* and y_1^* :
 - *graphical solution*: intersection of two lines
 - *numerically*: system of two linear equations

- for player 1:
$$\left. \begin{array}{l} EZ_1 = 11 - 10x_1 \\ EZ_1 = 2 + 5x_1 \end{array} \right\} \begin{array}{l} x_1 = 3/5 \\ EZ_1 = 5 \end{array}$$

- for player 2:
$$\left. \begin{array}{l} EZ_1 = 7 - 6y_1 \\ EZ_1 = 2 + 9y_1 \end{array} \right\} \begin{array}{l} y_1 = 1/3 \\ EZ_1 = 5 \end{array}$$

- equilibrium mixed strategies:

$$\mathbf{x}^* = \begin{bmatrix} \frac{3}{5} \\ \frac{2}{5} \end{bmatrix}, \quad \mathbf{y}^* = \begin{bmatrix} \frac{1}{3} \\ \frac{2}{3} \end{bmatrix}.$$

- value of the game: $\mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^* = 5$
 - represented by EZ_1 -value of the intersection in both plots

Graphical Solution: Exercise 2

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- *penalty kicks* game, kick better when aiming left
- use graphical solution to find NE strategies for both players

		Goalkeeper	
		Right	Left
Kick	1 \ 2		
	Right	30%	80%
Left	90%	50%	

Row and Column Differences Formula

- the solution of the system of two equations can be expressed easily using *row* and *column differences*:

$$r_1 = a_{11} - a_{12} \quad c_1 = a_{11} - a_{21}$$

$$r_2 = a_{21} - a_{22} \quad c_2 = a_{12} - a_{22}$$

- NE strategies:

$$x_1^* = \frac{r_2}{c_2 - c_1}, \quad y_1^* = \frac{c_2}{r_2 - r_1}$$

		Player 2		
		L	R	
Player 1	T	1	7	-6
	B	11	2	9
		-10	5	

a_{11}	a_{12}	r_1
a_{21}	a_{22}	r_2
c_1	c_2	

Row and Column Differences Formula

- derivation of the formula for NE strategies from the systems of equations:

$$\begin{array}{l}
 L: \quad EZ_1 = a_{11}x_1 + a_{21}(1 - x_1) = a_{21} + c_1x_1 \\
 R: \quad EZ_1 = a_{12}x_1 + a_{22}(1 - x_1) = a_{22} + c_2x_1
 \end{array}
 \left. \vphantom{\begin{array}{l} L \\ R \end{array}} \right\} x_1 = \frac{r_2}{c_2 - c_1}$$

$$\begin{array}{l}
 T: \quad EZ_1 = a_{11}y_1 + a_{12}(1 - y_1) = a_{12} + r_1y_1 \\
 B: \quad EZ_1 = a_{21}y_1 + a_{22}(1 - y_1) = a_{21} + r_2y_1
 \end{array}
 \left. \vphantom{\begin{array}{l} T \\ B \end{array}} \right\} y_1 = \frac{c_2}{r_2 - r_1}$$

		Player 2		
		L	R	
Player 1	T	1	7	-6
	B	11	2	9
		-10	5	

a_{11}	a_{12}	r_1
a_{21}	a_{22}	r_2
c_1	c_2	

Row and Column Differences: Exercise 3

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- using *row* and *column differences* formula to find the NE strategies in the following game

		Player 2	
		L	R
Player 1	T	3	7
	B	10	4

Graphical Solution – $2 \times n$ Matrices

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- a generalized version of 2×2 graphical solution
 - start with player 1's plot
 - determine the “active” best responses (strategies that can be played in a NE)
 - consider only active best responses for player 2's plot
- for $m \times 2$ matrices, proceed similarly

		Player 2		
		L	M	R
Player 1	T	1	12	7
	B	11	0	2

Expected payoffs with $\mathbf{x} = [x_1, 1-x_1]^T$:

$$L: \quad EZ_1 = 1 \times x_1 + 11 \times (1 - x_1) = 11 - 10x_1$$

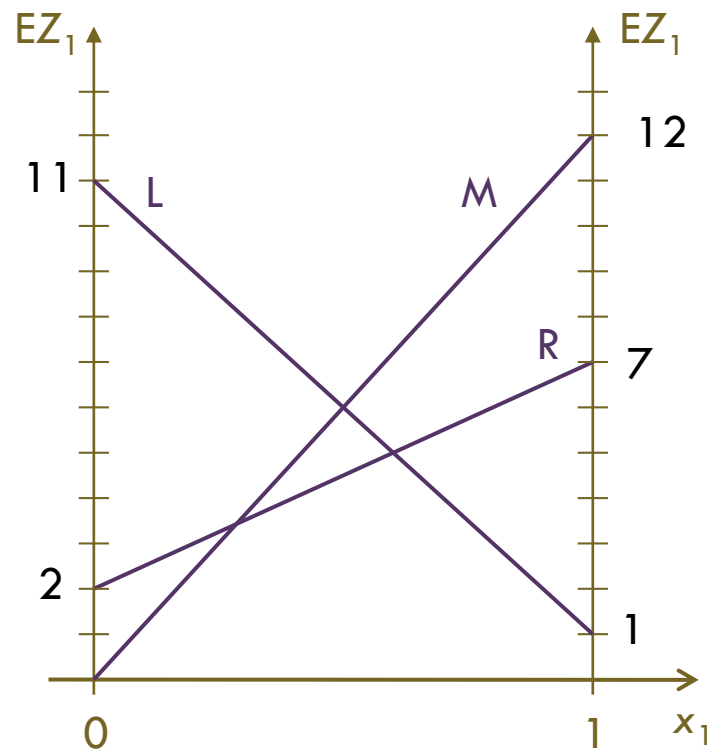
$$M: \quad EZ_1 = 12 \times x_1 + 0 \times (1 - x_1) = 12x_1$$

$$R: \quad EZ_1 = 7 \times x_1 + 2 \times (1 - x_1) = 2 + 5x_1$$

Graphical Solution – $2 \times n$ Matrices

(cont'd)

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Graphical Solution – $2 \times n$ Matrices

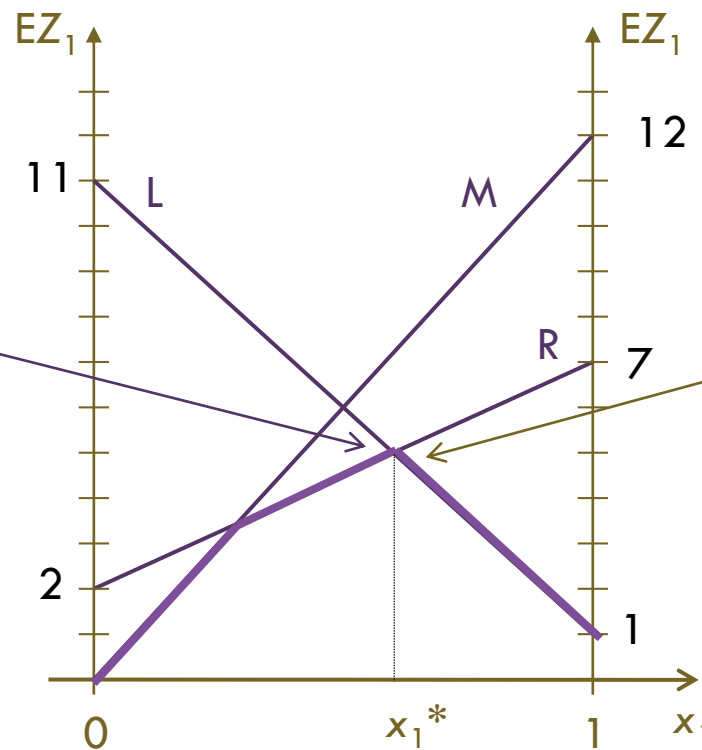
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active strategies of
player 2: L and R



$$y^* = \begin{bmatrix} y_1 & L \\ 0 & M \\ 1 - y_1 & R \end{bmatrix}$$

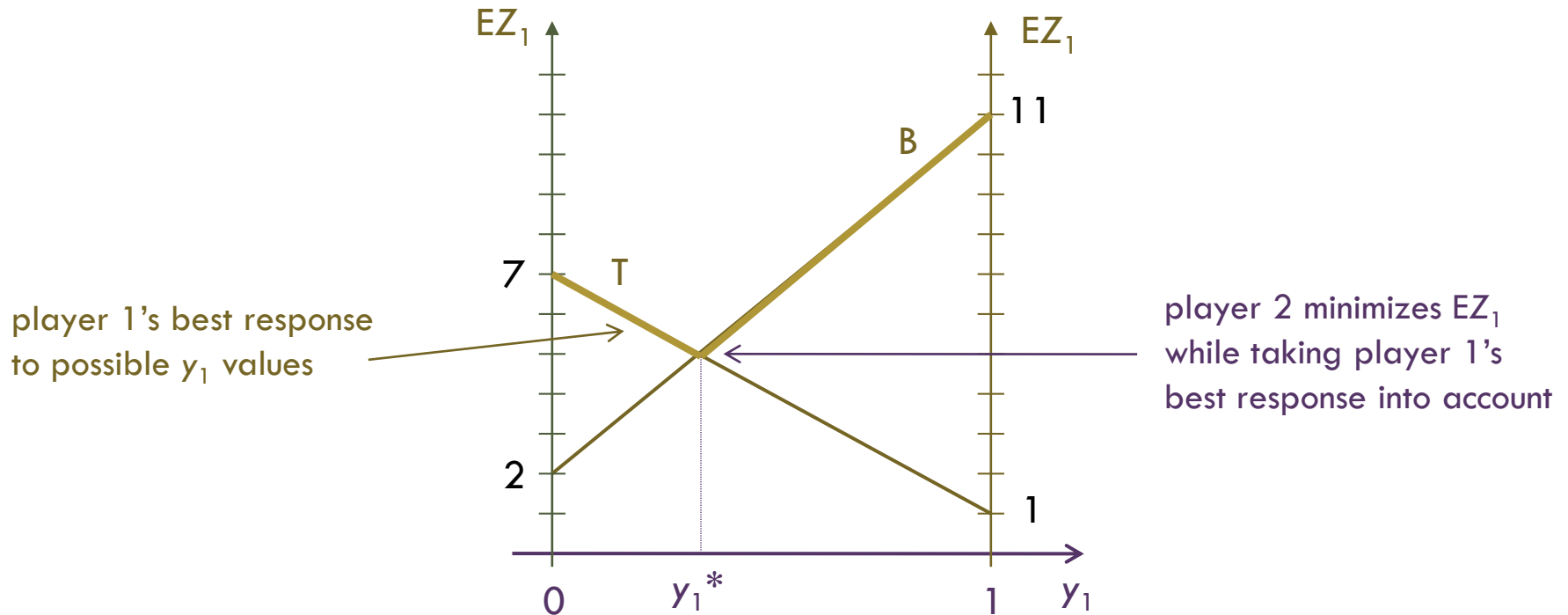


player 1 maximizes EZ_1
while taking player 2's
best response into account

Graphical Solution – $2 \times n$ Matrices

(cont'd)

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Finding NE – Linear Programming

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- the construction of the algorithm and its explanation is actually presented in the proof of the *Basic Theorem on Matrix Games* (see the *Games and Economic Decisions* textbook, or next lecture)

Linear programming:

Using *linear programming* methods, one can find a maximum or minimum of a linear function of multiple variables on a set given by linear constraints:

$$\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}, \mathbf{c} = \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_n \end{bmatrix}, \mathbf{A} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{bmatrix}, \mathbf{b} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{bmatrix} \quad \Rightarrow \quad \begin{array}{l} \text{maximize } z = \mathbf{c}^\top \mathbf{x} \\ \text{subject to} \\ \mathbf{Ax} \leq \mathbf{b}, \\ \mathbf{x} \geq \mathbf{0}. \end{array}$$

variables
data

- **Step 1:** If there is a negative element in the payoff matrix, make all elements of the matrix positive by adding the same positive number to all elements of the matrix. (This does changes the game, but only into a *strategically equivalent* one.)

- **Step 2:** Solve linear programming problem

maximize $p_1 + p_2 + \dots + p_n$

subject to

$$a_{11}p_1 + a_{12}p_2 + \dots + a_{1n}p_n \leq 1,$$

$$a_{21}p_1 + a_{22}p_2 + \dots + a_{2n}p_n \leq 1,$$

.....

$$a_{m1}p_1 + a_{m2}p_2 + \dots + a_{mn}p_n \leq 1,$$

$$p_i \geq 0, \quad i = 1, \dots, n.$$

- **Step 3:** Divide the primal and dual solutions by the optimal value of the objective function:
 - the *primal solution* determines the strategy of *player 2*.
 - the *dual solution* determines the strategy of *player 1*.

- note: if we use the symbol $\mathbf{1}_n$ to denote vector

$$\mathbf{1}_n = \left. \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \right\} n \text{ elements}$$

we can simplify the LP problem from step 2 as

$$\text{maximize } z = \mathbf{1}_n^\top \mathbf{p}$$

subject to

$$\mathbf{A}\mathbf{p} \leq \mathbf{1}_m,$$

$$\mathbf{p} \geq \mathbf{0}.$$

Using LP to Find NE: An Example

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- we'll find mixed-strategy NE's in the following matrix game:

		Player 2		
		L	M	R
Player 1	T	1	1.5	-1
	B	0	4	2

- **Step 1:** elimination of negative elements. We'll add a constant $c = 2$ to all elements of the matrix (to get a strategically equivalent matrix game with non-negative elements.)

		Player 2		
		L	M	R
Player 1	T	3	3.5	1
	B	2	6	4

Using LP to Find NE: An Example

(cont'd)

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- **Step 2:** solve LP problem

maximize $p_1 + p_2 + p_3$

subject to

$$3p_1 + 3.5p_2 + 1p_3 \leq 1,$$

$$2p_1 + 6p_2 + 4p_3 \leq 1,$$

$$p_1, p_2, p_3 \geq 0.$$

- solve using *MS Excel* (to see how, download the `NE_Solver.xls` file from my website), optimal values are:

- objective function value: 0.4

- primal solution: $\mathbf{p} = [0.3 \ 0 \ 0.1]^\top$

- dual solution: $\mathbf{q} = [0.2 \ 0.2]^\top$

- **Step 3:** the equilibrium strategies are:

- player 1: $\mathbf{x}^* = \mathbf{q} / 0.4 = [0.5 \ 0.5]^\top$

- player 2: $\mathbf{y}^* = \mathbf{p} / 0.4 = [0.25 \ 0 \ 0.75]^\top$

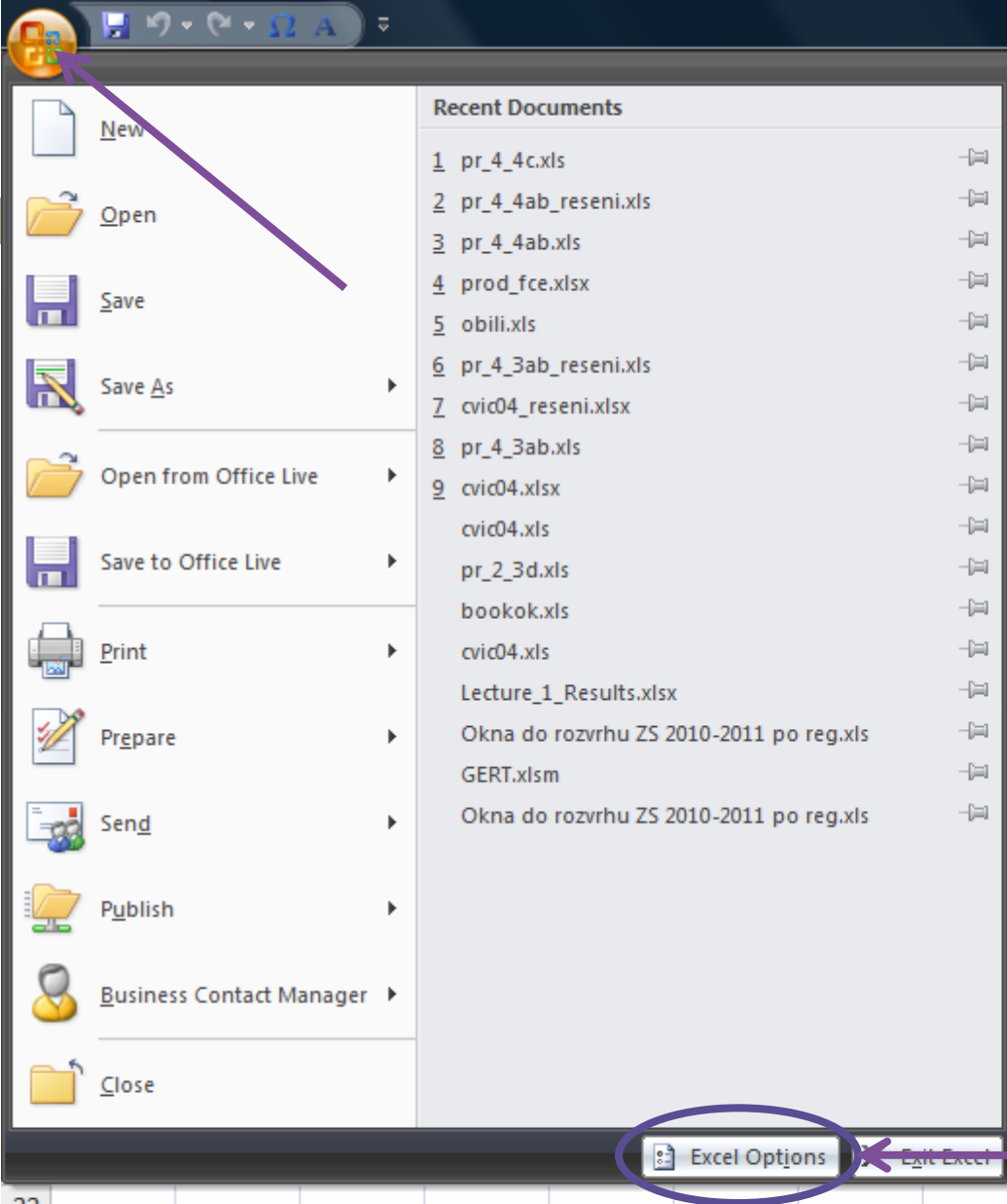
- value of the game is

$$\mathbf{x}^{*\top} \mathbf{A} \mathbf{y}^* = 1 / \text{objective function} = 1 / 0.4 = 2.5$$

- value of the original game is

$$1 / \text{objective function} - c = 1 / 0.4 - 2 = 2.5 - 2 = 0.5$$

following slides: activating the Solver add-in in MS Excel 2007



22
23
24
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26
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Excel Options

Popular

Formulas

Proofing

Save

Advanced

Customize

Add-Ins

Trust Center

Resources



View and manage Microsoft Office add-ins.

Add-ins

Name	Location	Type
Active Application Add-ins		
Business Contact Manager for Outlook	C:\...ess Contact Manager\BcmHistoryAddin.dll	COM Add-in
Microsoft Office Live Add-in	C:\...iles\Microsoft\Office Live\OLConnector.dll	COM Add-in
PDFComplete	C:\...m Files\PDF Complete\Addin\officepdf.dll	COM Add-in
Send to Bluetooth	C:\Windows\System32\btsendto_office.dll	COM Add-in
Inactive Application Add-ins		
Analytické nástroje	C:\...ce\Office12\Library\Analysis\ANALYS32.XLL	Excel Add-in
Analytické nástroje – VBA	C:\...Office12\Library\Analysis\ATPVBAEN.XLAM	Excel Add-in
Custom XML Data	C:\...es\Microsoft Office\Office12\OFFRHD.DLL	Document Inspector
Date (Smart tag lists)	C:\...iles\microsoft shared\Smart Tag\MOFL.DLL	Smart Tag
Financial Symbol (Smart tag lists)	C:\...iles\microsoft shared\Smart Tag\MOFL.DLL	Smart Tag
Headers and Footers	C:\...es\Microsoft Office\Office12\OFFRHD.DLL	Document Inspector
Hidden Rows and Columns	C:\...es\Microsoft Office\Office12\OFFRHD.DLL	Document Inspector
Hidden Worksheets	C:\...es\Microsoft Office\Office12\OFFRHD.DLL	Document Inspector
Internet Assistant – VBA	C:\...rosoft Office\Office12\Library\HTML.XLAM	Excel Add-in
Invisible Content	C:\...es\Microsoft Office\Office12\OFFRHD.DLL	Document Inspector
Nástroje pro měnu euro	C:\... Office\Office12\Library\EUROTOOL.XLAM	Excel Add-in
Person Name (Outlook e-mail recipients)	C:\...es\microsoft shared\Smart Tag\FNAME.DLL	Smart Tag
Průvodce podmíněným součtem	C:\...rosoft Office\Office12\Library\SUMIF.XLAM	Excel Add-in
Průvodce vyhledáváním	C:\...oft Office\Office12\Library\LOOKUP.XLAM	Excel Add-in
Document Related Add-ins		

Add-in: Business Contact Manager for Outlook

Publisher: Microsoft Corporation

Location: C:\Program Files\Microsoft Small Business\Business Contact Manager\BcmHistoryAddin.dll

Description: Communication History Addin

Manage: Excel Add-ins

Go...

OK

Cancel

LECTURE 3:
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